Computer Organization and Architecture Designing for Performance

11th Edition



Chapter 13

Instruction Sets: Characteristics and Functions



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Machine Instruction Characteristics

- The operation of the processor is determined by the instructions it executes, referred to as machine instructions or computer instructions
- The collection of different instructions that the processor can execute is referred to as the processor's instruction set
- Each instruction must contain the information required by the processor for execution

Figure 13.1 $\bigcirc \bigcirc \bigcirc$ **Instruction Cycle State Diagram**



Source and result operands can be in one of four areas:

- 1) Main or virtual memory
 - As with next instruction references, the main or virtual memory address must be supplied

2) I/O device

The instruction must specify the I/O module and device for the operation. If memorymapped I/O is used, this is just another main or virtual memory address

- 3) Processor register
 - A processor contains one or more registers that may be referenced by machine instructions.
 - If more than one register exists each register is assigned a unique name or number and the instruction must contain the number of the desired register

Immediate

The value of the operand is contained in a field in the instruction being executed

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Figure 13.2 O O O O O O O

- Within the computer each instruction is represented by a sequence of bits
- The instruction is divided into fields, corresponding to the constituent elements of the instruction



Instruction Representation

- Opcodes are represented by abbreviations
 called *mnemonics*
- Examples include:

ADD

SUB Subtract

Add

- MUL Multiply
- DIV Divide
- LOAD Load data from memory
 - STOR Store data to memory
- Operands are also represented symbolically
- Each symbolic opcode has a fixed binary representation
 - The programmer specifies the location of each symbolic operand

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Instruction Types



Figure 13.3 O O O O O O O Figure 13.3 Figu

\bigcirc	Instruction SUBY, A, B MPYTDE	$\frac{Comment}{Y \leftarrow A - B}$ $T \leftarrow D \times E$				-0
0	ADD T, T, C DIV Y, Y, T	$T \leftarrow T + C$ $Y \leftarrow Y \div T$				-0
\bigcirc		8	J	Instruction	Comment	-0
\bigcirc	(a) Three-	address instructions		LOAD D	AC ← D	-0
	.	<i>a i</i>		MPY E	$AC \leftarrow AC \times E$	
\bigcirc	MOVE Y A	Comment V ← A		ADD C	$AC \leftarrow AC + C$	
	SUB Y, B	$Y \leftarrow Y - B$		LOAD A	$AC \leftarrow A$	
	MOVE T, D	T←D		SUB B	$AC \leftarrow AC - B$	
	MPY T, E	$T \leftarrow T \times E$	h	DIV Y	$AC \leftarrow AC \div Y$	
	ADD I.C	$I \leftarrow I + C$ $V \leftarrow V \div T$		STOR Y	$Y \leftarrow AC$	()
	1,1	6 6		5 6	6	\bigcirc
						8

Table 13.1Image: Construction of the second sec

Number of Addresses	Symbolic Representation	Interpretation
3	OP A, B, C	A ← B OP C
2	OP A, B	A ← A OP B
1	OP A	$AC \leftarrow AC OP A$
0	OP	T ← (T – 1) OP T
AC T T (T-1) = A, B, C =	accumulator top of stack second element of stack memory or register locations	





Numbers

- All machine languages include numeric data types
- Numbers stored in a computer are limited:
 - Limit to the magnitude of numbers representable on a machine
 - In the case of floating-point numbers, a limit to their precision
- Three types of numerical data are common in computers:
 Binary integer or binary fixed point
 - Binary floating point
 - Decimal
- Packed decimal
 - Each decimal digit is represented by a 4-bit code with two digits stored per byte
 - To form numbers 4-bit codes are strung together, usually in multiples of 8 bits

Characters

- A common form of data is text or character strings
- Textual data in character form cannot be easily stored or transmitted by data processing and communications systems because they are designed for binary data
- Most commonly used character code is the International Reference Alphabet (IRA)
 - Referred to in the United States as the American Standard Code for Information Interchange (ASCII)
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- Another code used to encode characters is the Extended Binary Coded Decimal Interchange Code (EBCDIC)
 - EBCDIC is used on IBM mainframes

Logical Data

- An *n*-bit unit consisting of *n* 1-bit items of data, each item having the value 0 or 1
- Two advantages to bit-oriented view:
 - Memory can be used most efficiently for storing an array of Boolean or binary data items in which each item can take on only the values 1 (true) and 0 (false)
 - To manipulate the bits of a data item
 - If floating-point operations are implemented in software, we need to be able to shift significant bits in some operations
 - To convert from IRA to packed decimal, we need to extract the rightmost 4 bits of each byte

Data Type	Description	Table 13 2
General	Byte, word (16 bits), doubleword (32 bits), quadword (64 bits), and double quadword (128 bits) locations with arbitrary binary contents.	OC Data Trace
Integer	A signed binary value contained in a byte, word, or doubleword, using twos complement representation.	xoo Data Types
Ordinal	An unsigned integer contained in a byte, word, or doubleword.	
Unpacked binary coded decimal (BCD)	A representation of a BCD digit in the range 0 through 9, with one digit in each byte.	
Packed BCD	Packed byte representation of two BCD digits; value in the range 0 to 99.	– –
Near pointer	A 16-bit, 32-bit, or 64-bit effective address that represents the offset within a segment. Used for all pointers in a nonsegmented memory and for references within a segment in a segmented memory.	
Far pointer	A logical address consisting of a 16-bit segment selector and an offset of 16, 32, or 64 bits. Far pointers are used for memory references in a segmented memory model where the identity of a segment being accessed must be specified explicitly.	
Bit field	A contiguous sequence of bits in which the position of each bit is considered as an independent unit. A bit string can begin at any bit position of any byte and can contain up to 32 bits.	E
Bit string	A contiguous sequence of bits, containing from zero to $2^{23} - 1$ bits.	22
Byte string	A contiguous sequence of bytes, words, or doublewords, containing from zero to $2^{23} - 1$ bytes.	
Floating point	See Figure 13.4.	
Packed SIMD (single instruction, multiple data)	Packed 64-bit and 128-bit data types.	(Table can be found on page 443

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Single-Instruction-Multiple-Data (SIMD) Data Types

- Introduced to the x86 architecture as part of the extensions of the instruction set to optimize performance of multimedia applications
- These extensions include MMX (multimedia extensions) and SSE (streaming SIMD extensions)
- Data types:
 - Packed byte and packed byte integer
 - Packed word and packed word integer
 - Packed doubleword and packed doubleword integer
 - Packed quadword and packed quadword integer
 - Packed single-precision floating-point and packed doubleprecision floating-point





(a) Data Transfer

Operation Name	Description
MOV Dest, Source	Move data between registers or between register and memory or immediate to register.
XCHG Op1, Op2	Swap contents between two registers or register and memory.
PUSH Source	Decrements stack pointer (ESP register), then copies the source operand to the top of stack.
POP Dest	Copies top of stack to destination and increments ESP.

Table 13.3 Common x86 Instruction

(b) Arithmetic

Operation Name	Description		
ADD Dest, Source	Adds the destination and the source operand and stores the result in the destination. Destination can be register or memory. Source can be register, memory, or immediate.		
SUB Dest, Source	Subtracts the source from the destination and stores the result in the destination.		
MUL Op	Unsigned integer multiplication of the operand by the AL, AX, or EAX register and stores in the register. Opcode indicates size of register.		
IMUL Op	Signed integer multiplication.		
DIV Op	Divides unsigned the value in the AX, DX:AX, EDX:EAX, or RDX:RAX registers (dividend) by the source operand (divisor) and stores the result in the AX (AH:AL), DX:AX, EDX:EAX, or RDX:RAX registers.		
IDIV Op	Signed integer division.		
INC Op	Adds 1 to the destination operand, while preserving the state of the CF flag.		
DEC Op	Subtracts 1 from the destination operand, while preserving the state of the CF flag.		
NEG Op	Replaces the value of operand with (0 – operand), using twos complement (Ta representation.		
CMP Op1, Op2	Compares the two operands by subtracting the second operand from the first operand and sets the status flags in the EFLAGS register according to the results.		

Set Operations (1 of 3)

Table can be found on page 446-447 in the textbook.)

(c) Shift and Rotate

Operation Name	Description
SAL Op, Quantity	Shifts the source operand left by from 1 to 31 bit positions. Empty bit positions are cleared. The CF flag is loaded with the last bit shifted out of the operand.
SAR Op, Quantity	Shifts the source operand right by from 1 to 31 bit positions. Empty bit positions are cleared if the operand is positive and set if the operand is negative. The CF flag is loaded with the last bit shifted out of the operand.
SHR Op, Quantity	Shifts the source operand right by from 1 to 31 bit positions. Empty bit positions are cleared and the CF flag is loaded with the last bit shifted out of the operand.
ROL Op, Quantity	Rotate bits to the left, with wraparound. The CF flag is loaded with the last bit shifted out of the operand.
ROR Op, Quantity	Rotate bits to the right, with wraparound. The CF flag is loaded with the last bit shifted out of the operand.
RCL Op, Quantity	Rotate bits to the left, including the CF flag, with wraparound. This instruction treats the CF flag as a one-bit extension on the upper end of the operand.
RCR Op, Quantity	Rotate bits to the right, including the CF flag, with wraparound. This instruction treats the CF flag as a one-bit extension on the lower end of the operand.

Table 13.3 Common x86 Instruction Set

Operations

(2 of 3)

Operation Name	Description			
NOT Op	Inverts each bit of the operand.			
AND Dest, Source	Performs a bitwise AND operation on the destination and source operands and stores the result in the destination operand.			
OR Dest, Source	Performs a bitwise OR operation on the destination and source operands and stores the result in the destination operand.			
XOR Dest, Source	Performs a bitwise XOR operation on the destination and source operands and stores the result in the destination operand.			
TEST Op1, Op2	Performs a bitwise AND operation on the two operands and sets the S, Z, and P status flags. The operands are unchanged.			

(d) Logical

(Table can be found on page 446-447 in the textbook.)

Universidad de Costa Rica 21

(e) Transfer of Control

Operation Name	Description
CALL proc	Saves procedure linking information on the stack and branches to the called procedure specified using the operand. The operand specifies the address of the first instruction in the called procedure.
RET	Transfers program control to a return address located on the top of the stack. The return is made to the instruction that follows the CALL instruction.
JMP Dest	Transfers program control to a different point in the instruction stream without recording return information. The operand specifies the address of the instruction being jumped to.
Jcc Dest	Checks the state of one or more of the status flags in the EFLAGS register (CF, OF, PF, SF, and ZF) and, if the flags are in the specified state (condition), performs a jump to the target instruction specified by the destination operand. See Tables 13.8 and 13.9.
NOP	This instruction performs no operation. It is a one-byte or multi-byte NOP that takes up space in the instruction stream but does not impact machine context, except for the EIP register.
HLT	Stops instruction execution and places the processor in a HALT state. An enabled interrupt, a debug exception, the BINIT# signal, the INIT# signal, or the RESET# signal will resume execution.
WAIT	Causes the processor to repeatedly check for and handle pending, unmasked, floating- point exceptions before proceeding.
INT Nr	Interrupts current program, runs specified interrupt program

(f) Input/Output

Operation Name	Description
IN Dest, Source	Copies the data from the I/O port specified by the source operand to the destination operand, which is a register location.
INS Dest, Source	Copies the data from the I/O port specified by the source operand to the destination operand, which is a memory location.
OUT Dest, Source	Copies the byte, word, or doubleword value from the source register to the I/O port specified by the destination operand.
OUTS Dest, Source	Copies byte, word, or doubleword from the source operand to the I/O port specified with the destination operand. The source operand is a memory location

Table 13.3 Common x86 Instruction Set

Operations

(3 of 3)

(Table can be found on page 446-447 in the textbook.)

Table 13.4Image: Constraint of the second secon

		Transfer data from one location to another	
0—	Data transfer	If memory is involved: Determine memory address	
0—		Perform virtual-to-actual-memory address transformation Check cache Initiate memory read/write	0
		May involve data transfer, before and/or after	
	Arithmetic	Perform function in ALU	
		Set condition codes and flags	
\bigcirc	Logical	Same as arithmetic	
	Conversion	Similar to arithmetic and logical. May involve special logic to perform conversion	
	Transfer of control	Update program counter. For subroutine call/return, manage parameter passing and linkage	
		Issue command to I/O module	
	"0	If memory-mapped I/O, determine memory-mapped address	

(Table can be found on page 447 in the textbook.)

Data Transfer

Most fundamental type of machine instruction Must specify:

Location of the

- source and destination operands
- The length of data to be transferred must be indicated
- The mode of addressing for each operand

must be specified

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Table 13.5Image: Constraint of the second secon

Operation Mnemonic	Name	Number of Bits Transferred	Description	
L	Load	32	Transfer from memory to register	
LH	Load Halfword	16	Transfer from memory to register	
LR	Load	32	Transfer from register to register	
LER	Load (short)	32 Transfer from floating-point register to floating-point register		
LE	Load (short)	32	Transfer from memory to floating-point register	
LDR	Load (long)	64	Transfer from floating-point register to floating-point register	
LD	Load (long)	64	Transfer from memory to floating-point register	
ST	Store	32	Transfer from register to memory	
STH	Store Halfword	16	Transfer from register to memory	
STC	Store Character	8	Transfer from register to memory	
STE	Store (short)	32	Transfer from floating-point register to memory	
STD	Store (long)	64	Transfer from floating-point register to memory	

Arithmetic

- Most machines provide the basic arithmetic operations of add, subtract, multiply, and divide
- These are provided for signed integer (fixed-point) numbers
- Often they are also provided for floating-point and packed decimal numbers
- Other possible operations include a variety of single-operand instructions:
 - Absolute
 - Take the absolute value of the operand
 - Negate
 - Negate the operand
 - Increment
 - Add 1 to the operand
 - Decrement
 - Subtract 1 from the operand

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Table 13.6Image: Constraint of the second secon

Р	Q	NOT P	P AND Q	P OR Q	P XOR Q	P = Q
0	0	1	0	0	0	1
0	1	1	0	1	1	0
1	0	0	0	1	1	0
1	1	0	1	1	0	1
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Table 13.7Image: Constraint of the second secon

Input	Operation	Result	
10100110	Logical right shift (3 bits)	00010100	
10100110	Logical left shift (3 bits)	00110000	
10100110	Arithmetic right shift (3 bits)	11110100	
10100110	Arithmetic left shift (3 bits)	10110000	
10100110	Right rotate (3 bits)	11010100	
10100110	Left rotate (3 bits)	00110101	

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Instructions that change the format or operate on the format of data

> An example of a more complex editing instruction is the EAS/390 Translate (TR) instruction

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An example is converting from decimal to binary

6

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Conversion

Input/Output

DMA

- Variety of approaches taken:
 - Isolated programmed I/O
 - Memory-mapped programmed I/O
 - Use of an I/O processor
- Many implementations provide only a few I/O instructions, with the specific actions specified by parameters, codes, or command words

System Control

Instructions that can be executed only while the processor is in a certain privileged state or is executing a program in a special privileged area of memory

Typically these instructions are reserved for the use of the operating system

Examples of system control operations:

A system control instruction may read or alter a control register An instruction to read or modify a storage protection key Access to process control blocks in a multiprogramming system

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Transfer of Control

- Reasons why transfer-of-control operations are required:
 - It is essential to be able to execute each instruction more than once
 - Virtually all programs involve some decision making
 - It helps if there are mechanisms for breaking the task up into smaller pieces that can be worked on one at a time
- Most common transfer-of-control operations found in instruction sets:
 - Branch
 - Skip
 - Procedure call

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Figure 13.7 **O O O O O O**



Skip Instructions

Includes an implied address

Typically implies that one instruction be skipped, thus the implied address equals the address of the next instruction plus one instruction length

Because the skip instruction does not require a destination address field it is free to do other things



Example is the increment-and-skip-ifzero (ISZ) instruction

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Procedure Call Instructions

- Self-contained computer program that is incorporated into a larger program
 - At any point in the program the procedure may be invoked, or *called*
 - Processor is instructed to go and execute the entire procedure and then return to the point from which the call took place
- Two principal reasons for use of procedures:

Economy

• A procedure allows the same piece of code to be used many times

- Modularity
- Involves two basic instructions:
 - A call instruction that branches from the present location to the procedure
 - Return instruction that returns from the procedure to the place from which it was called







x86 Operation Types

- The x86 provides a complex array of operation types including a number of specialized instructions
- The intent was to provide tools for the compiler writer to produce optimized machine language translation of high-level language programs
- Provides four instructions to support procedure call/return:
 - CALL
 - ENTER
 - LEAVE
 - RETURN
- When a new procedure is called the following must be performed upon entry to the new procedure:
 - Push the return point on the stack
 - Push the current frame pointer on the stack
 - Copy the stack pointer as the new value of the frame pointer
 - Adjust the stack pointer to allocate a frame

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Table 13.8Image: Constraint of the second secon

Status Bit	Name	Description	
С	Carry	Indicates carrying or borrowing out of the left- most bit position following an arithmetic operation. Also modified by some of the shift and rotate operations.	
Р	Parity	Parity of the least-significant byte of the result of an arithmetic or logic operation. 1 indicates even parity; 0 indicates odd parity.	
А	Auxiliary Carry	Represents carrying or borrowing between half- bytes of an 8-bit arithmetic or logic operation. Used in binary- coded decimal arithmetic.	
Z	Zero	Indicates that the result of an arithmetic or logic operation is 0.	
S	Sign	Indicates the sign of the result of an arithmetic or logic operation.	
Ο	Overflow	Indicates an arithmetic overflow after an addition or subtraction for twos complement arithmetic.	
S O	Sign Overflow	Indicates the sign of the result of an arithmetic or logic operation. Indicates an arithmetic overflow after an addition or subtraction for twos complement arithmetic.	

Symbol	Condition Tested	Comment
A, NBE	C = 0 AND Z = 0	Above; Not below or equal (greater than, unsigned)
AE, NB, NC	C = 0	Above or equal; Not below (greater than or equal, unsigned); Not carry
B, NAE, C	C = 1	Below; Not above or equal (less than, unsigned); Carry set
BE, NA	C = 1 OR Z = 1	Below or equal; Not above (less than or equal, unsigned)
E, Z	Z = 1	Equal; Zero (signed or unsigned)
G, NLE	[(S = 1 AND O = 1) OR (S = 0 AND O = 0)]AND[Z = 0]	Greater than; Not less than or equal (signed)
GE, NL	(S = 1 AND O = 1) OR (S = 0 AND O = 0)	Greater than or equal; Not less than (signed)
L, NGE	(S = 1 AND O = 0) OR (S = 0 AND O = 0)	Less than; Not greater than or equal (signed)
LE, NG	(S = 1 AND O = 0) OR (S = 0 AND O = 1) OR (Z = 1)	Less than or equal; Not greater than (signed)
NE, NZ	Z = 0	Not equal; Not zero (signed or unsigned)
NO	O = 0	No overflow
NS	S = 0	Not sign (not negative)
NP, PO	P = 0	Not parity; Parity odd
0	O = 1	Overflow
Р	P = 1	Parity; Parity even
S	S = 1	Sign (negative)

Table 13.9 x86 Condition **Codes for** Conditional **Jump and SETcc Instructions** 022

(Table can be found on page 460 in the textbook.) 42

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x86 Single-Instruction, Multiple-Data (SIMD) Instructions

- 1996 Intel introduced MMX technology into its Pentium product line
 - MMX is a set of highly optimized instructions for multimedia tasks
- Video and audio data are typically composed of large arrays of small data types

- Three new data types are defined in MMX
 - Packed byte
 - Packed word
 - Packed doubleword
- Each data type is 64 bits in length and consists of multiple smaller data fields, each of which holds a fixed-point integer

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MMX Instruction Set	Arithn
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Category	Instruction	Description	
	PADD [B, W, D]	Parallel add of packed eight bytes, four 16-bit words, or two 32-bit doublewords, with wraparound.	
	PADDS [B, W]	Add with saturation.	
	PADDUS [B, W]	Add unsigned with saturation.	
	PSUB [B, W, D]	Subtract with wraparound.	
	PSUBS [B, W]	Subtract with saturation.	
Arithmetic	PSUBUS [B, W]	Subtract unsigned with saturation.	
	PMULHW	Parallel multiply of four signed 16-bit words, with high- order 16 bits of 32-bit result chosen.	
	PMULLW	Parallel multiply of four signed 16-bit words, with low-order 16 bits of 32-bit result chosen.	
	PMADDWD	Parallel multiply of four signed 16-bit words; add together adjacent pairs of 32-bit results.	
Comparison	PCMPEQ [B, W, D]	Parallel compare for equality; result is mask of 1s if true or 0s if false.	
	PCMPGT [B, W, D]	Parallel compare for greater than; result is mask of 1s if true or 0s if false.	
	PACKUSWB	Pack words into bytes with unsigned saturation.	
	PACKSS [WB, DW]	Pack words into bytes, or doublewords into words, with signed saturation.	
Conversion	PUNPCKH [BW, WD, DQ]	Parallel unpack (interleaved merge) high- order bytes, words, or doublewords from MMX register.	
	PUNPCKL [BW, WD, DQ]	Parallel unpack (interleaved merge) low- order bytes, words, or doublewords from MMX register.	
	PAND	64-bit bitwise logical AND	
Logical	PNDN	64-bit bitwise logical AND NOT	
	POR	64-bit bitwise logical OR	
	PXOR	64-bit bitwise logical XOR	
Shift	PSLL [W, D, Q]	Parallel logical left shift of packed words, doublewords, or quadword by amount specified in MMX register or immediate value.	
	PSRL [W, D, Q]	Parallel logical right shift of packed words, doublewords, or quadword.	
	PSRA [W, D]	Parallel arithmetic right shift of packed words, doublewords, or quadword.	
Data transfer	MOV [D, Q]	Move doubleword or quadword to/from MMX register.	
Statemgt	EMMS	Empty MMX state (empty FP registers tag bits).	

Note: If an instruction supports multiple data types [byte (B), word (W), doubleword (D), quadword (Q)], the data types are indicated in brackets. (Table can be found on page 462 in the textbook.) (Table can be found on page 462 in the textbook.)



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ARM Operation Types

Load and store instructions

Branch instructions

Multiply instructions

Parallel addition and subtraction instructions

Extend instructions

46

Data-

processing

instructions

Status register access instructions

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	Code	S
Table 13.11	0000	EQ
ARM Conditions	0001	NE
for Conditional	0010	CS/
	0011	CC/
Instruction	00100	MI
Execution	00101	ΡL
Execution	00110	VS
	00111	VC
	1000	н
	1001	LS
	1010	GE
	1011	LT
	1100	GT
	1101	LE
	1110	AL
2	1111	-

Code	Symbol	Condition Tested	Comment
0000	EQ	Z = 1	Equal
0001	NE	Z = 0	Not equal
0010	CS/HS	C = 1	Carry set/unsigned higher or same
0011	CC/LO	C = 0	Carry clear/unsigned lower
00100	МІ	N = 1	Minus/negative
00101	PL	N = 0	Plus/positive or zero
00110	VS	V = 1	Overflow
00111	VC	V = 0	No overflow
1000	н	C = 1 AND Z = 0	Unsigned higher
1001	LS	C = 0 OR Z = 1	Unsigned lower or same
1010	GE	N = V [(N = 1 AND V = 1) OR (N = 0 AND V = 0)]	Signed greater than or equal
1011	LT	N ≠ V [(N = 1 AND V = 0) OR (N = 0 AND V = 1)]	Signed less than
1100	GT	(Z = 0) AND (N = V)	Signed greater than
1101	LE	(Z = 1) OR (N ≠ V)	Signed less than or equal
1110	AL	_	Always (unconditional)
1111	-	-	This instruction can only be executed unconditionally

(Table can be found on page 465 in the textbook.)

Summary

Chapter 13

- Machine instruction characteristics
 - Elements of a machine
 instruction
 - Instruction representation
 - Instruction types
 - Number of addresses
 - Instruction set design
- Types of operands
 - Numbers
 - Characters
 - Logical data

Intel x86 and ARM data types

Instruction Sets:

Characteristics and

Functions

- Types of operations
 - Data transfer
 - Arithmetic
 - Logical
 - Conversion
 - Input/output
 - System control
 - Transfer of control
 - Intel x86 and ARM operation types